Building Highly Reliable Networks with GRASP/VND Heuristics

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Motivation

Remark

- Reliability falls within the field of metrology.
- Its practical interests is found in network design.
- Choose links to maximize reliability & connectivity.
- Minimum-cost topologies are partially known (eg. kECON).
- What happens under probabilistic models?

Main Goal

We are given $2r$ terminals.
Choose $3r$ links to achieve maximum reliability.
Definition (Unreliability)

The unreliability of a simple graph $G$ with independent link failures with probability $\rho$ is:

$$U_G(\rho) = \sum_{k=0}^{q} m_k \rho^k (1 - \rho)^{q-k},$$

being $m_k$ the number of ways to disconnect $G$ removing $k$ links. A $(\rho, q)$-graph is a graph with $p$ nodes and $q$ links.

Definition (Uniformly Most-Reliable Graph (UMR))

A $(\rho, q)$-graph $G$ is UMR if its unreliability is minimum among all $(\rho, q)$-graphs and all $\rho \in [0, 1]$. 
Motivation

Problem

Research Process

Results

Conclusions

Research Process

Necessary and Sufficient Conditions

Proposition (Sufficient Condition)

If $m_k(G) \leq m_k(H)$ for all $k$ and all $(p, q)$-graph $H$, then $G$ is uniformly most-reliable.

Proposition (Necessary Condition)

Optimal graphs $G$ must have the maximum tree-number $\tau(G)$, maximum connectivity $\lambda(G)$, and the minimum number of disconnecting sets $m_\lambda(G)$. 
Known Cubic UMR Graphs

(a) $K_4 = M_2$  (b) $K_{(3,3)} = M_3$  (c) $M_4$

Figure: Complete, Bipartite and Wagner graphs.

(d) $P_5$  (e) $Y_6$

Figure: Petersen and Yutsis graphs.
Minimize each coefficient $m_k$ in individual blocks.
Combine them in a Variable Neighborhood Search (VND).
The previous idea is computationally prohibitive!
At least, force the necessary conditions.
Use GRASP/VND to return a candidate.
The objective should be only one...
Conjecture (Boesch et. al.)

If $G$ is uniformly most-reliable $(p, q)$-graph, then $m_k(G) \leq m_k(H)$ for all $(p, q)$-graph $H$.

Remark

If Boesch conjecture holds, then the number of disconnected subgraphs $m(G) = \sum_{k=0}^{q} m_k$ must be minimized.

Finding $m(G)$ is $\mathcal{NP}$-Hard (an evaluation of Tutte polynomial). Observe that $m(G) = 2^q \times U_G(1/2)$. An estimation for $U_G(1/2)$ is available using Monte Carlo.
Algorithm 1  \( G = \text{HighlyReliable}(r, \text{iter}, \alpha) \)

1: \( G \leftarrow M_r \)
2: \textbf{for} \( i = 1 \) \textbf{to} \( \text{iter} \) \textbf{do}
3: \( G_{\text{input}} \leftarrow \text{Construct}(r, \alpha) \)
4: \( G(i) \leftarrow \text{VND}(G_{\text{input}}) \)
5: \textbf{if} \( m(G(i)) \leq m(G) \) \textbf{for all} \( k \) \textbf{then}
6: \( G \leftarrow G(i) \)
7: \textbf{end if}
8: \textbf{end for}
9: \textbf{return} \( G \)
Research Process

VND - Flow Diagram

\[ G \leftarrow \text{Construct}(r, \alpha) \]

Is \( G \) Healthy?

\[ \text{no} \]

\[ G \leftarrow \text{Surgery}(G) \]

Is \( G \) Cubic?

\[ \text{no} \]

\[ G \leftarrow \text{Cubic}(G) \]

Is \( G \) Strong?

\[ \text{yes} \]

\[ G \leftarrow \text{Crossing}(G) \]

\[ \text{Return } G \]
Experimental Setting

Settings

- \( r \in \{7, \ldots, 15\} \).
- \( \text{iter} = 10^6 \) (executions of the GRASP/VND)
- \( N = 10^4 \) (sample graphs in Monte Carlo)
- Output: 9 graphs: \( R_r \).
- Brute force test: they are the only UMR candidates.
Results: $R_7$ and $R_8$

Heawood Graph

$U_{R_7}(\frac{1}{2}) = 0.9177$
$\tau(R_7) = 50421$

Möbius-Kantor Graph

$U_{R_8}(\frac{1}{2}) = 0.94320$
$\tau(R_8) = 248832$
Results: $R_9$ and $R_{10}$

$U_{R_9}(\frac{1}{2}) = 0.9607$
$\tau(R_9) = 1265625$

$U_{R_{10}}(\frac{1}{2}) = 0.97310$
$\tau(R_{10}) = 6422000$
Results: $R_{11}$ and $R_{12}$

$U_{R_{11}}\left(\frac{1}{2}\right) = 0.9801$
$\tau(R_{11}) = 32710656$

$U_{R_{12}}\left(\frac{1}{2}\right) = 0.988$
$\tau(R_{12}) = 168664320$
**Results:** $R_{13}$ and $R_{14}$

\[
U_{R_{13}}\left(\frac{1}{2}\right) = 0.9904 \\
\tau(R_{13}) = 862488000
\]

\[
U_{R_{14}}\left(\frac{1}{2}\right) = 0.9931 \\
\tau(R_{14}) = 4410450000
\]
Results: $R_{15}$

\[ U_{R_{15}} \left( \frac{1}{2} \right) = 0.9956 \]
\[ \tau(R_{15}) = 23066015625 \]
Conclusions

1. Finding UMR graphs is hard.
2. A methodology to find highly reliable networks is proposed.
3. New candidates of UMR graphs are also found.
4. UMR graphs are highly symmetrical.
5. Several conjectures are still open.
6. Only few works deal with node-failures.